

# A Hash Function Family *Luffa*

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# Outline -

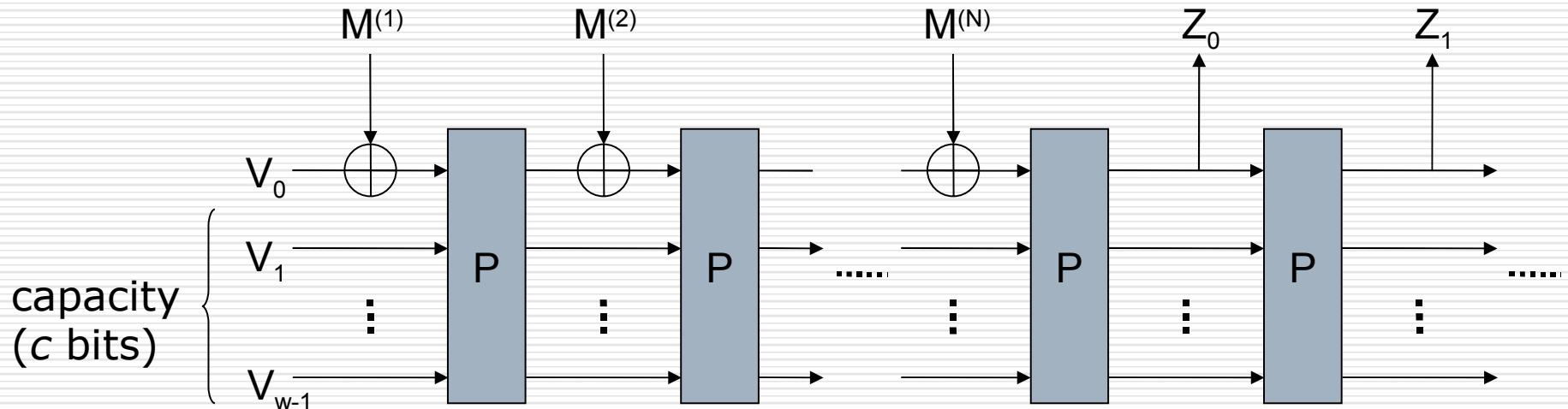
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- Specification
  - Chaining -
  - Non-linear components
- Security status
  - Generic attack
  - Differential based attack
- Implementations
  - Software
  - Hardware

# Introduction to *Luffa* (spec.)

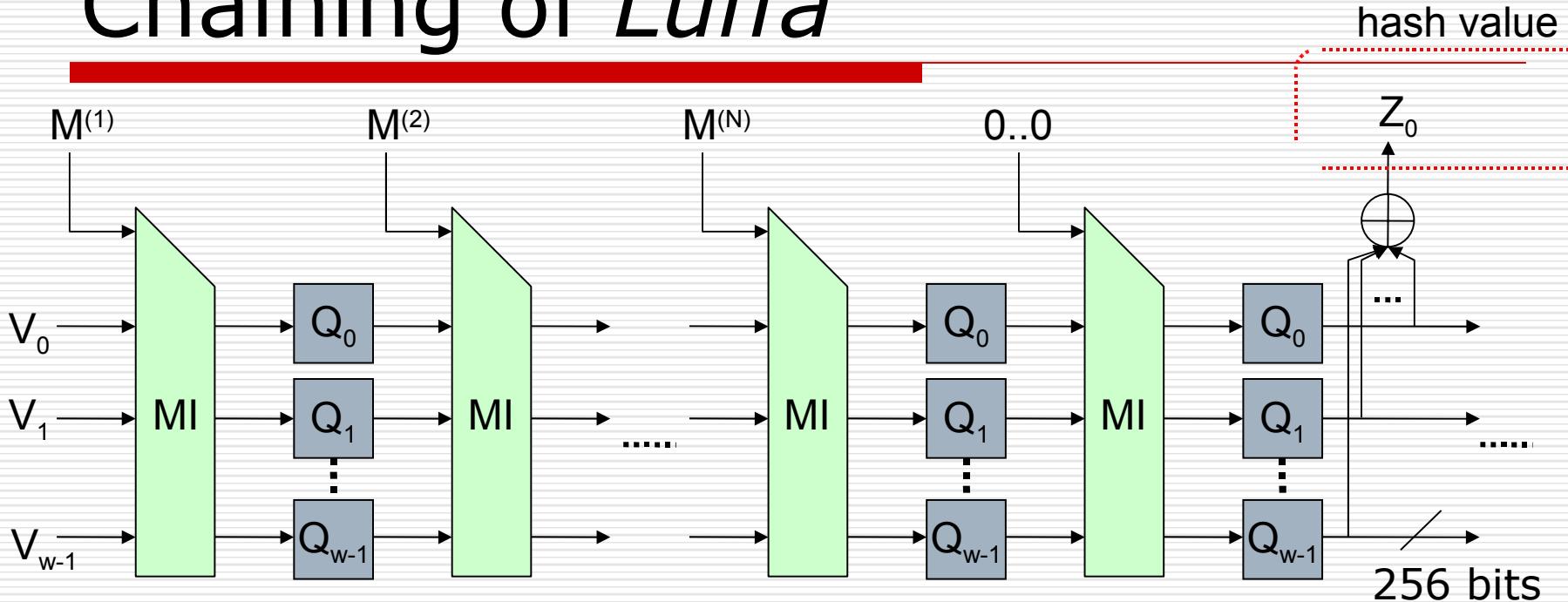
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# Cryptographic sponge function



- Newer coming construction of a hash function from a random permutation
- It is indifferentiable from a RO

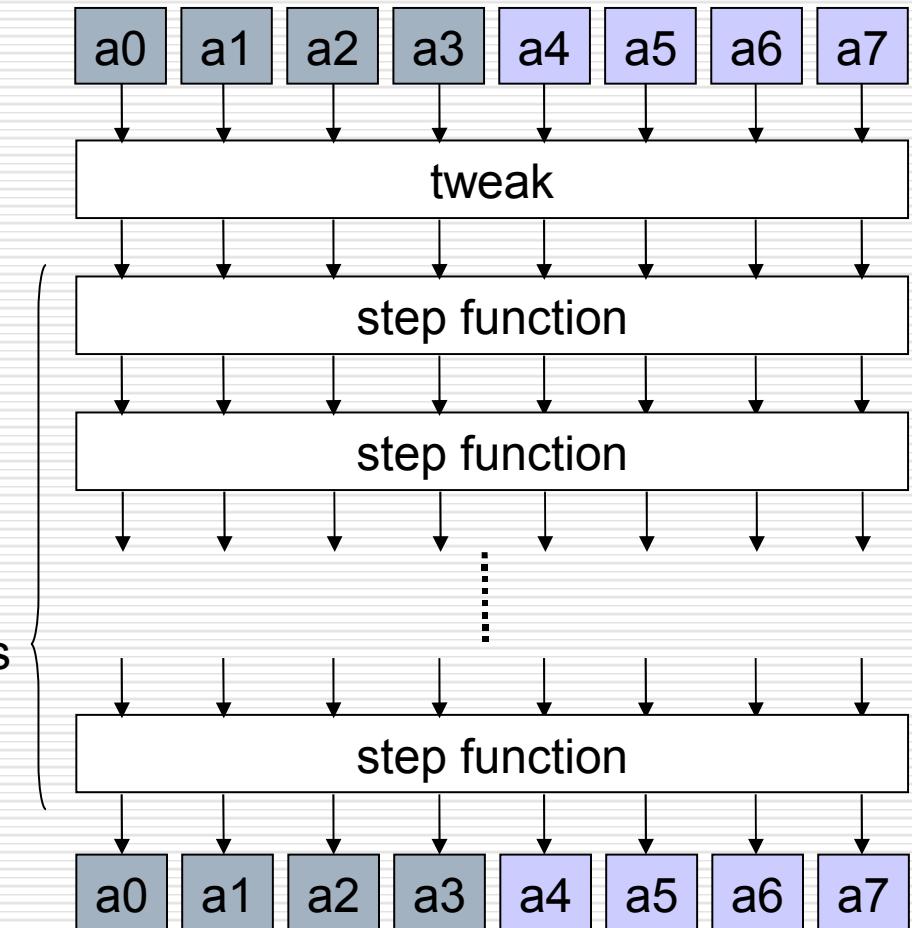
# Chaining of *Luffa*



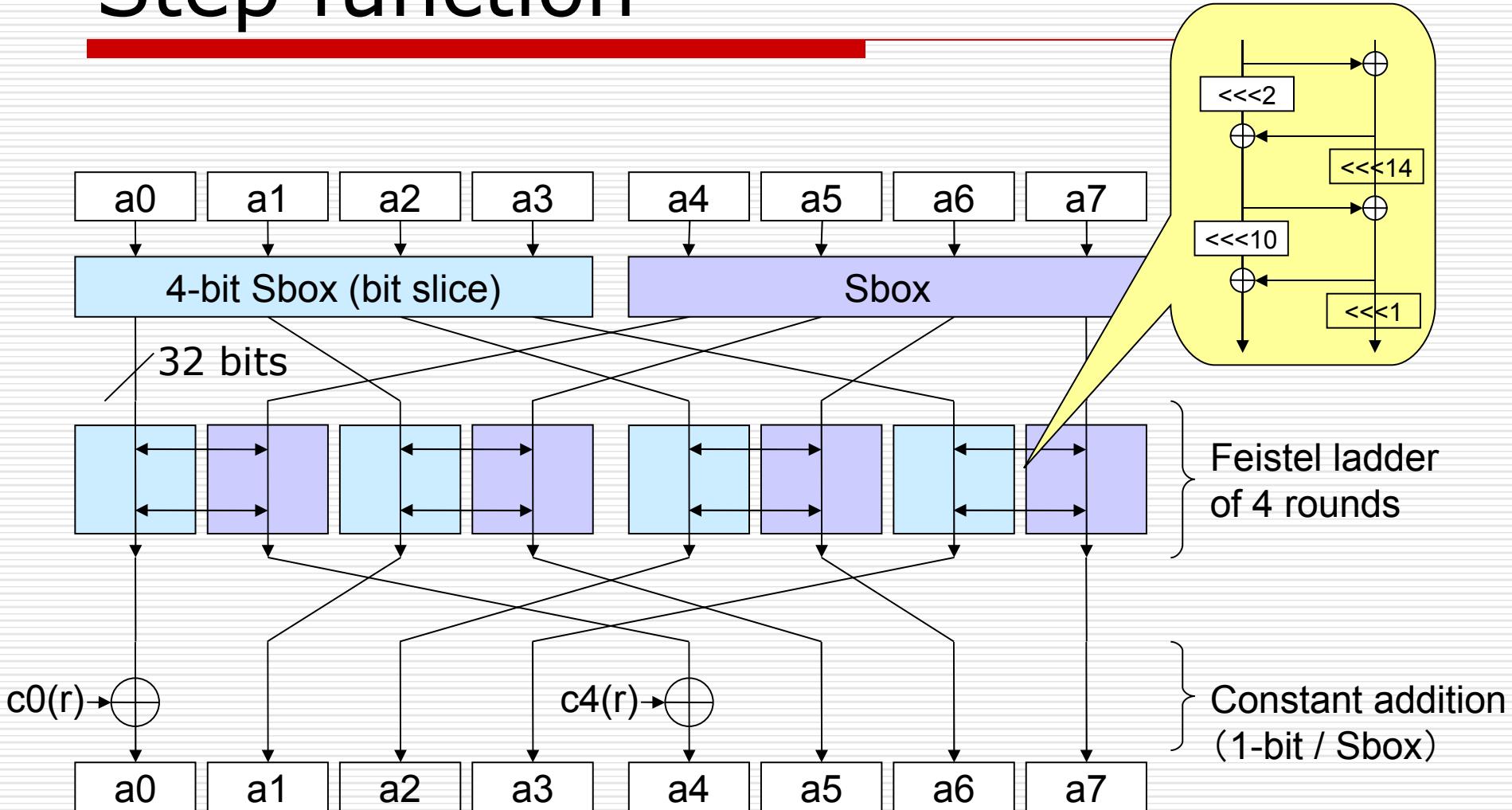
- *Luffa* is a variant of sponge
  - But, fixed length permutations for all hash length -
    - The number of  $Q_j$  increases if the hash length gets long ( $w=3, 4, 5$  for  $\text{hash\_len}=256, 384, 512$ )
  - Insert message and mix the state by the linear map  $MI$
  - A blank round
  - The hash value is the sum of the outputs of  $Q_j$

# Non-linear permutation Q -

- Input/Output
  - 256 bits  
(8 32-bit words)
- Functions
  - tweak -
    - Applied before step functions -
  - Step functions -
    - 8 steps



# Step function



# Security status

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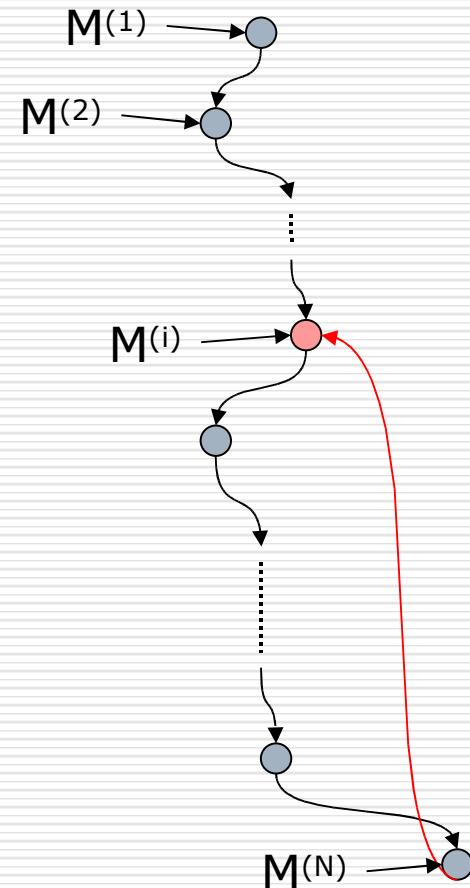
# Summary of security status -

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- Sponge function features
  - Not based on CR compression function
  - Finding inner collision is the best attack
- Current security status of *Luffa*
  - No security proof for the chaining (yet)
  - Several generic attacks concerned, none of them are serious
  - Differential based attack
    - Seems secure under a reasonable assumption

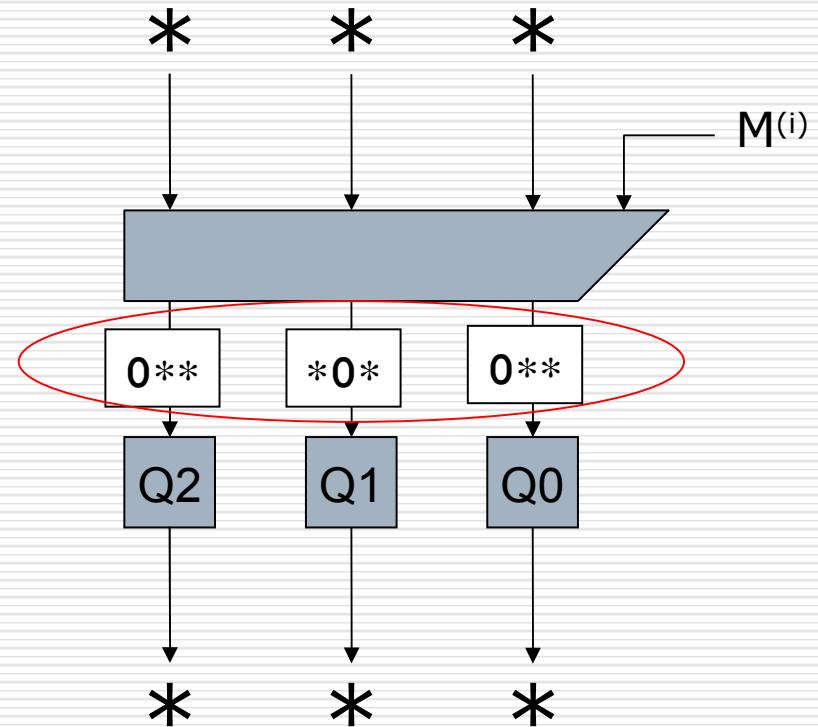
# Long message attack -

- Sponge's case
  - Finding a message s.t.  
 $S^{(i)} = S^{(N)}$
  - Prob. of the event
    - capacity:  
 $c = \text{len}(S) - \text{len}(M)$
    - prob =  $2^{-c/2}$
  - Complexity -
    - Queries to the permutation:  $2^{c/2}$
    - Num. of nodes:  $2^{c/2}$



# Long message attack (conti.)

- Luffa's case
  - 1/w of input bits to each  $Q_j$  is controllable by message injection
- Complexity
  - Queries to  $Q_j$ 
    - $2^{(w-1)/w} \cdot 256$
  - Num. of nodes
    - $2^{(w-1)/2} \cdot 256$
  - Calc. Complexity
    - MA:  $2^{(w-1)/2} \cdot 256$
    - MI calls:  $2^{(w-1)/2} \cdot 256$



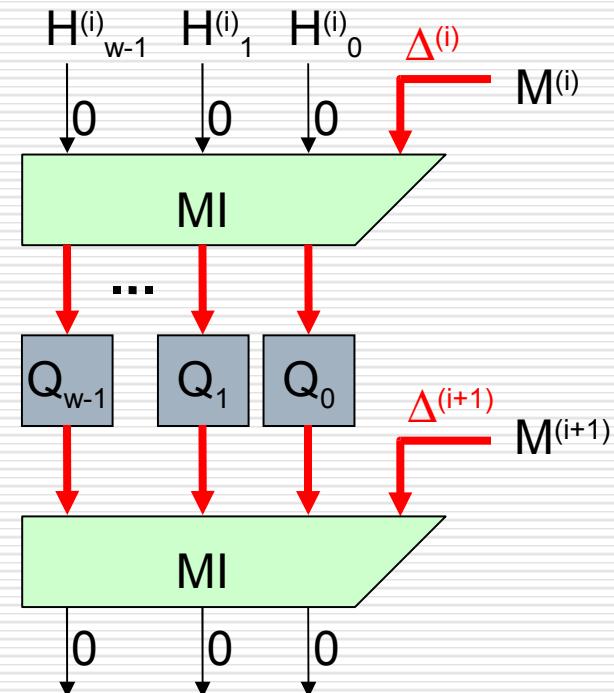
# Differential characteristics of Qj-

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- 4 steps (half-block)
  - Approach: exhaustive truncated path search -
  - Possible min. num. of active Sbox: 31
  - MDCP  $\leq 2^{-62}$
- 8 steps (full)
  - Approach: Leon's algorithm to find the lowest code word
  - Min. active Sbox = 112
  - DCP =  $2^{-224} (> 2^{-256})$
  - Not useful to find an inner collision

# Differential based attack scenario

- (Seems) the best scenario
  - 2 rounds attack to find an - inner collision -
- Limitation of modification technique
  - Assumption -
    - 1 bit modification doubles - the diff. prob. -
  - Message block  $M^{(i)}$ 
    - Any, up to 256 bits
  - State  $H^{(i)}$ -
    - Assumed random, up to  $(w-1)/2 \cdot 256$  bits -
- (Our) conclusion
  - *Luffa* is secure against this attack if  $MDCP(Q_j) < 2^{-171}$



# Implementation aspects -

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# Software implementations -

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hash length	ANSI C (cycle/byte)		assembly with SSE2 (cycle/byte)	
	32-bit	64-bit	32-bit	64-bit
224	33.9	32.0	13.9	13.4
256	33.4	32.0	13.9	13.4
384	45.2	39.0	15.7	15.2
512	59.7	50.3	25.5	23.2

- Evaluation environment
  - CPU: Intel Core2Duo E6600 (2.4GHz)
  - Memory: 2GB
  - ANSI C: Windows Vista + Visual Studio 2005
  - Assembly: Ubuntu Linux 8.04 + gas

# Hardware implementations (ASIC)

Hash length (bit)	Opt.	Gate count (gate)	Frequency (MHz)	Cycles	Throughput (Mbps)
256/224	size	10,157	100	891	28.7
256/224	speed	26,849	444	9	12,642
384	speed	34,985	444	9	12,642
512	speed	44,163	444	9	12,642

- Evaluation environment
  - 0.13µm CMOS standard cell library
- Optimization
  - Small gate size: with 1 Sbox and 1 MixWord
  - Speed: 3 step functions in parallel



# Thanks you for your attention! -

# FAQ -

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- Q1. What is *Luffa*?
  - A vegetable sponge
  - Scientific name: *Luffa cylindrica* (See picture) -
- Q2. Why *Luffa*?
  - Because it is a kind of - sponge -
  - And very useful (like as hash function) -
    - High-quality sponge from dried fruit -
    - The young fruit is edible
    - Face lotion from the juice
    - Educational material (in Japan) -
    - This is the first trial to use - *Luffa* in cryptography -



Photo reprinted from Wikipedia