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# The HIGHT Encryption Algorithm draft-kisa-hight-00

#### Abstract

This document describes the HIGHT(HIGH security and light weigHT) encryption algorithm, which is suitable for low-resource device. HIGHT is a 64-bit block cipher with 128-bit keys. The algorithm consists of round functions, key schedule, encryption, and decryption.

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#### 1 Introduction

## 1.1. HIGHT overview

HIGHT is a 64-bit symmetric key light-weight block cipher suitable for low-resource device. HIGHT stands for 'HIGh security and light weigHT' and is developed by Korea 2005. HIGHT is a ISO/IEC international standard block cipher algorithm included in ISO/IEC 18033-3:2010 [ISO-HIGHT]. It has simple structure with use of basic arithmetic operation - XOR, addition/subtraction in modular 2\*\*8, and circular shift rotation, and also without S-Box.

The features of HIGHT are outlined as follows:

- 64-bit input/output data block size
- 128-bit key length
- 32-round with XOR, modular addition, and shift rotation
- No S-Box
- Designed for low-resource device (data storage, power, etc.)

# 2. Notation and Terminology

The following notation is used in the description of HIGHT encryption algorithm:

[^] bitwise XOR addition in modular 2\*\*8 [+] [-] subtraction in modular 2\*\*8  $\Pi$ concatenation left circular shift rotation by n-bit in 8-bit value <<<n Р plaintext ciphertext С master key K WK whitening key SK subkey round function 0 F0 F1 round function 1 i-th byte of X Χi i-th byte of  $\boldsymbol{X}$  in round j Xj,i di intermediate status value in subkey scheduling

#### 3. The HIGHT algorithm

## 3.1. Round functions

The HIGHT algorithm uses two round functions, namely, F0 and F1 which are now defined.

a) Round function 0 (F0) The F0 function is used for encryption and decryption with 8-bit input. The function F0 is defined as follows:

$$F0(x) = (x <<<1) [^] (x <<<2) [^] (x <<<7)$$

b) Round function 1 (F1) The F1 function is used for encryption and decryption with 8-bit input. The function F1 is defined as follows:

$$F1(x) = (x <<<3) [^] (x <<<4) [^] (x <<<6)$$

## 3.2. Key schedule

The key schedule for HIGHT describes the procedure to make whitening key bytes WKi and 128 subkey bytes SKi from a 128-bit master key K = K15 || K14 || ... || K0, as shown below.

a) The generation of whitening keys is defined as follows.

for 
$$i = 0, 1, 2, 3$$
:  
 $WKi = K(i+12)$   
for  $i = 4, 5, 6, 7$ :  
 $WKi = K(i-4)$ 

- b) The 128 subkeys are used for encryption and decryption, 4 subkeys per round. The generation of subkeys is defined as follows.
  - (1) s0 = 0, s1 = 1, s2 = 0, s3 = 1, s4 = 1, s5 = 0, s6 = 1d0 = s6 || s5 || s4 || s3 || s2 || s1 || s0
  - (2) for i = 1 to 127:  $s(i+6) = s(i+2) \lceil \wedge \rceil s(i-1)$ di = s(i+6)||s(i+5)||s(i+4)||s(i+3)||s(i+2)||s(i+1)||si
  - (3) for i = 0 to 7: for j = 0 to 7:  $SK(16*i+j) = K(j-i \mod 8) [+] d(16*i+j)$ for j = 0 to 7:  $SK(16*i+j+8) = K((j-i \mod 8)+8)$  [+] d(16\*i+j+8)

#### 3.3. HIGHT encryption

The encryption operation is as shown in Figure 1. The transformation of a 64-bit block P into a 64-bit block C is defined as follows

```
(1) P = P7 || P6 || P5 || P4 || P3 || P2 || P1 || P0
(2) X0,0 = P0 [+] WK0,
                                    X0,1 = P1,
    X0,2 = P2 [^] WK1,
                               X0,3 = P3,
X0,5 = P5,
    X0,4 = P4 [+] WK2,
                                  X0,7 = P7.
    X0,6 = P6 [^] WK3,
(3) for i = 0 to 30:
      X(i+1), 0 = Xi, 7 [^] (FO(Xi, 6)[+]SK(4*i+3)), X(i+1), 1 = Xi, 0,
      X(i+1),2 = Xi,1 [+] (F1(Xi,0)[^]SK(4*i)), X(i+1),3 = Xi,2,
      X(i+1), 4 = Xi, 3 [^] (F0(Xi, 2)[+]SK(4*i+1)), X(i+1), 5 = Xi, 4,
      X(i+1), 6 = Xi, 5 [+] (F1(Xi, 4)[^]SK(4*i+2)), X(i+1), 7 = Xi, 6.
    for i = 31:
      X(i+1),0 = Xi,0, X(i+1),1 = Xi,1 [+] (F1(Xi,0)[^]SK124), X(i+1),2 = Xi,2, X(i+1),3 = Xi,3 [^] (F0(Xi,2)[+]SK125), X(i+1),4 = Xi,4, X(i+1),5 = Xi,5 [+] (F1(Xi,4)[^]SK126), X(i+1),6 = Xi,6, X(i+1),7 = Xi,7 [^] (F0(Xi,6)[+]SK127).
(4) C0 = X32,0 [+] WK4,
                                   C1 = X32, 1,
    (5) C = C7 || C6 || C5 || C4 || C3 || C2 || C1 || C0
```

#### 3.4. HIGHT Decryption

The decryption operation is identical in operation to encryption apart from the following two modifications.

- (1) All [+] operations are replaced by [-] operations except for the [+] operations connecting SKi and outputs of F0.
- (2) The order in which the keys WKi and SKi are applied is reversed.

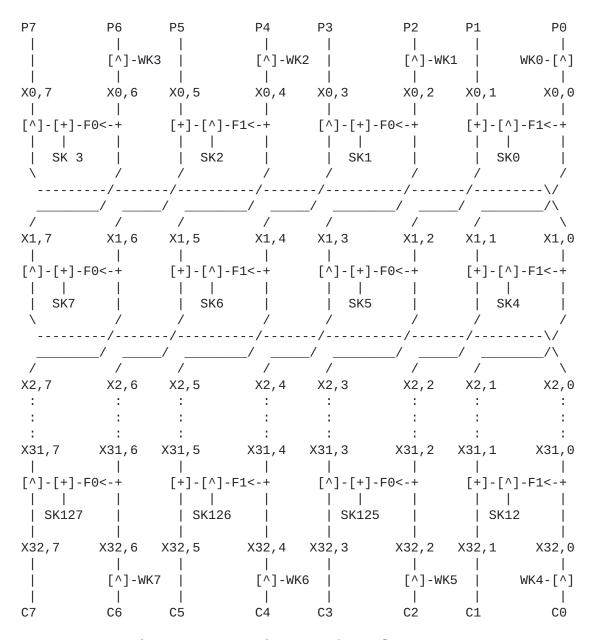


Figure 1. Encryption procedure of HIGHT

## 3.5. HIGHT Object Identifiers

For those who may be using HIGHT in algorithm negotiation within a protocol, or in any other context that may require the use of Object Identifiers (OIDs), the following OIDs have been defined.

```
algorithm OBJECT IDENTIFIER ::= { iso(1) member-body(2) korea(410)
          kisa(200004) algorithm(1) }
id-hight OBJECT IDENTIFIER ::= { algorithm hight(40) }
id-hightECB OBJECT IDENTIFIER ::= { algorithm hightECB(41) }
id-hightCBC OBJECT IDENTIFIER ::= { algorithm hightCBC(42) }
id-hightCFB OBJECT IDENTIFIER ::= { algorithm hightCFB(43) }
id-hightOFB OBJECT IDENTIFIER ::= { algorithm hightOFB(44) }
id-hightCTR OBJECT IDENTIFIER ::= { algorithm hightCTR(45) }
```

The id-hightECB, id-hightCBC, id-hightCFB, id-hightOFB, and idhightCTR OIDs are used when Electronic CodeBook (ECB) mode, Cipher Block Chaining (CBC) mode, Cipher Feed-Back (CFB) mode, Output Feed-Back (OFB) mode, and Counter (CTR) mode of operation based on the HIGHT block cipher is provided respectively.

#### 4. Security Considerations

No security problem has been found on HIGHT.

# <u>5</u>. Test Vectors

## 5.1. Test vectors 1

Key:
00 11 22 33 44 55 66 77 88 99 aa bb cc dd ee ff
Plaintext:
00 00 00 00 00 00 00 00
Ciphertext:
00 f4 18 ae d9 4f 03 f2

Sub Key	Value	Sub Key	Value
SK3   SK2   SK1   SK0	= e7135b59,	 SK67   SK66   SK65   SK64	= cfa7c7f6
SK7   SK6   SK5   SK4	= c99cb0c8,	SK71   SK70   SK69   SK68	= 48555f62
SK11  SK10  SK9   SK8	= 906d96d7,	SK75   SK74   SK73   SK72	= 1f50a1b1
SK15  SK14  SK13  SK12	= 2c6a5599,	SK79   SK78   SK77   SK76	= a5986d86
SK19  SK18  SK17  SK16	= 27032ade,	SK83   SK82   SK81   SK80	= 0706f33c
SK23  SK22  SK21  SK20	= b5e32d31,	SK87   SK86   SK85   SK84	= fb2b7aff
SK27  SK26  SK25  SK24	= ced9de4e,	SK91   SK90   SK89   SK88	= 7a755a93
SK31  SK30  SK29  SK28	= 48919180,	SK95   SK94   SK93   SK92	= 7bb39134
SK35  SK34  SK33  SK32	= f915b5f4,	SK99   SK98   SK97   SK96	= bcdf15f0
SK39  SK38  SK37  SK36	= fadc0ee2,	SK103  SK102  SK101  SK100	= ef018ca2
SK43  SK42  SK41  SK40	= bba15439,	SK107  SK106  SK105  SK104	= 2a436495
SK47  SK46  SK45  SK44	= 9fadb9bf,	SK111  SK110  SK109  SK108	= ae882255
SK51  SK50  SK49  SK48	= 16b7f8e8,	SK115  SK114  SK113  SK112	= c7e50f52
SK55  SK54  SK53  SK52	= e41e0239,	SK119  SK118  SK117  SK116	= 67d9bcf0
SK59  SK58  SK57  SK56	= d9451b36,	SK123  SK122  SK121  SK120	= 61a18fda
SK63  SK62  SK61  SK60	= a9b0ad97,	SK127  SK126  SK125  SK124	= d1357c79
=======================================			========

Round	Value	Round	Value
Initial	0000001100220033	 Round 17	2c93a90ddd0283ae
Round 1	00ce1138223f33e7	Round 18	93570db102d9aec4
Round 2	cee138ef3fa3e78a	Round 19	57b7b1dbd998c4e4
Round 3	e14fef91a3708a8a	Round 20	b7bedb55989ae458
Round 4	4f8a91cd70518ad1	Round 21	be87559d9a515868
Round 5	8a53cd0951c3d1ee	Round 22	87ce9d5351786873
Round 6	534609c7c3e4ee7d	Round 23	ceab53d6784b73bc
Round 7	4673c7c5e41b7dd7	Round 24	ab30d6d74ba8bc69
Round 8	7359c58c1b33d79c	Round 25	30bfd7f7a83369df
Round 9	595f8cf333d59c07	Round 26	bf13f71733bfdf7d
Round 10	5f0cf317d507073f	Round 27	134617f1bfd57db2
Round 11	0ca0173007033fb6	Round 28	467bf187d5c4b277
Round 12	a03a3043030bb63e	Round 29	7b3187d2c4f5772b
Round 13	3a7943b40b2b3e37	Round 30	315dd246f5482bde
Round 14	7920b47a2b7c37b5	Round 31	5d3846d148a1def3
Round 15	20637a797ce4b5d0	Round 32	003818d1d9a103f3
Round 16	632c79a9e4ddd083	Final	00f418aed94f03f2
========	:===========	========	=======================================

## 5.2. Test vectors 2

Key: ff ee dd cc bb aa 99 88 77 66 55 44 33 22 11 00 Plaintext: 00 11 22 33 44 55 66 77 Ciphertext: 23 ce 9f 72 e5 43 e6 d8

Sub Key	Value	Sub Key	Value
SK3   SK2   SK1   SK0   SK7   SK6   SK5   SK4   SK11  SK10  SK9   SK8   SK11  SK10  SK13  SK12   SK19  SK18  SK17  SK16   SK23  SK22  SK21  SK20   SK27  SK26  SK25  SK24   SK31  SK30  SK29  SK28   SK35  SK34  SK33  SK32   SK39  SK38  SK37  SK36   SK43  SK42  SK41  SK40	= 4e587e5a, = b8695b51, = 07c2c9e8, = 2b471032, = 6c262bcd, = 828eb698, = 230cef4d, = 254c2af7, = 1c16a4c1, = a5657527, = eeb25316,	SK67   SK66   SK65   SK64 SK71   SK70   SK69   SK68 SK75   SK74   SK73   SK72 SK79   SK78   SK77   SK76 SK83   SK82   SK81   SK80	= be74727f = af9a8263 = 1e2d5c4a = 1ceda097 = d4b17ca3 = 404e7bee = 5730f30a = d0e6a233 = 67687c35 = 12027b6f = e5dcdbea
SK51  SK50  SK49  SK48 SK55  SK54  SK53  SK52	= 17a6c593, = 6d85475c,	SK115  SK114  SK113  SK112 SK119  SK118  SK117  SK116	= 504c5475 = 68c8899b
		SK123  SK122  SK121  SK120 SK127  SK126  SK125  SK124	

Round		Value	Round	Value
Initia	 1	 00ee222144886643	 Round 17	db63ca6b6e9dfaaf
Round	1	ee2d21b1880a435f	Round 18	63776b6b9d09af72
Round	2	2db4b11c0acc5fde	Round 19	77856b93091172c5
Round	3	b4951c9fcca3dec5	Round 20	851793871106c58c
Round	4	95c19fe4a30fc556	Round 21	17a7878206f18c48
Round	5	c115e4730f545645	Round 22	a7598251f1c64855
Round	6	15e27386540d45b7	Round 23	597d5119c6e85575
Round	7	e26486c30dabb777	Round 24	7d4a196ee8e775d8
Round	8	6424c35bab9d7772	Round 25	4a7f6ef7e7bdd882
Round	9	24725b8c9d607282	Round 26	7fadf729bdcb8284
Round	10	72458c7b602d829d	Round 27	ad442985cb29845f
Round	11	458c7bab2dc69d59	Round 28	44b58548296e5f31
Round	12	8cc6ab08c6ba5982	Round 29	b51d488f6e0231f3
Round	13	c60f0841ba688280	Round 30	1df78ff802f8f39d
Round	14	0fd3413668f280d4	Round 31	f7fdf850f8529dd8
Round	15	d35c3627f2afd4e4	Round 32	23fd9f50e552e6d8
Round	16	5cdb27caaf6ee4fa	Final	23ce9f72e543e6d8

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#### 6. References

## **6.1.** Normative References

[ISO-HIGHT] ISO/IEC, "Information technology - Security techniques -Encryption - Part 3: Block ciphers", ISO/IEC 18033-3, December 2010.

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